



Software Analyzers

TEACHING FRAMA-C FOR CYBERSECURITY

Julien Signoles julien.signoles@cea.fr

23th May 2026 @ FMTea'26

CEA-List, Université Paris-Saclay, Software Safety and Security Lab

funded by the French National Research Agency through Project CMA-TCE (Train Cyber Experts)



list

université PARIS-SACLAY



anr agence nationale de la recherche



- 1 } Context
- 2 } Lecture and Lab Session
- 3 } Feedback

- > single **3-hour lecture** and single **3-hour lab session**
- > **shorter version** available:
 - > 90-minute lecture and 90-minute lab session
 - > 1- or 2-hour lecture, without lab session
- > created Q4 2022 for Telecom SudParis (engineering school), since:
 - > **Telecom SudParis** (engineering school), 2022-2025
 - > **CentraleSupélec Rennes** (engineering school), 2023-2025
 - > **Université Bretagne Sud** (international master students), 2023-2025
 - > **École des Mines Paris** (engineering school), 2025

> target audience:

- > students at master level
- > embedded system or cybersecurity track
- > no knowledge in formal methods needed

> requirement: reading C code (bachelor level)

- > knowledge of integer and pointer arithmetic, dynamic allocation, etc

> goals:

- > introduce formal methods to students
- > emphasize practical aspects in cybersecurity
- > show what they are good, but do not hide practical difficulties

Framework for analyzing C source code

<https://frama-c.com>

- > developed at CEA List since 2005: >20-year old
- > open source (LGPL 2.1)
- > based on formal methods
 - > ACSL as formal specification language
- > several code analysis tools provided as plug-ins
 - > >30 plug-ins in the latest distribution
- > plug-ins linked to the Frama-C kernel
 - > uniform setting (shared user interface, etc)
 - > information gathering (which analyzer proves which properties, etc)
 - > analysis collaboration

- 1 } Context
- 2 } Lecture and Lab Session
- 3 } Feedback

<https://julien-signoles.fr/teaching/slides-frama-c-cyber.pdf>

- 1 › Context
- 2 › Frama-C in a Nutshell
- 3 › Deductive Verification with ACSL and Wp
- 4 › Value Analysis with Eva
- 5 › Runtime Annotation Checking with E-ACSL
- 6 › Advanced Security Verification



breadth-based approach, instead of depth-based

- > why formal methods are important:
 - > for **trusting software**, particularly in critical systems

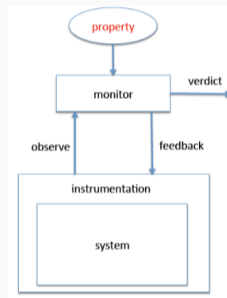
- > why it is not easy: **Rice theorem**
 - a **sound**, **automatic**, **exact** and **sound** analysis is not possible (undecidability)

- > how to be pragmatic: **different formal techniques**, none being perfect
 - > runtime verification (not static)
 - > model checking (does not always scale)
 - > abstract interpretation (not exact)
 - > deductive verification (not automatic)

Similar Outline for every technique presented:

- 1 › general overview of the technique: no maths!
- 2 › examples of basic usage and demo on open-source code
 - › live session on my laptop
- 3 › show practical difficulties and limitations
- 4 › industrial example

- > check whether a system under scrutiny satisfies a given property
- > **offline** monitoring: after execution
- > **online** monitoring: during execution
 - > **inline**: the monitor is part of the system
 - > **outline**: the monitor is outside the system
- > most approaches focus on verifying **temporal properties**
- > ... yet not E-ACSL



- > E-ACSL is a **Frama-C plug-in** that **generates inline monitors** for (E-)ACSL properties on C programs [Signoles et al, 2017]
 - > takes as input an (E-)ACSL-annotated C program
 - > generates a new C program
 - > that behaves as the original C program if all the annotations are valid; or
 - > fails on the first invalid annotation (by default)

<pre> 1 int div(int x, int y) { 2 /*@ assert y-1 != 0; */ 3 return x / (y-1); 4 }</pre>	<p>E-ACSL</p> <hr style="border: 0.5px solid orange; width: 100%;"/> <p>→</p>	<pre> 1 int div(int x, int y) { 2 /*@ assert y-1 != 0; */ 3 e_acsl_assert(y-1 != 0L); 4 return x / (y-1); 5 }</pre>
---	---	---



PolarSSL

(now known as Mbed-TLS)

<https://git.trustedfirmware.org/mirror/mbed-tls.git/about/>

- > C implementation of TLS (aka SSL)
- > not as complex as openSSL

SHOWCASING DIFFICULTIES: WHAT DO YOU (NOT) PROVE WITH WP?

- > what do you prove?
 - > the source code matches the spec
 - > under the assumptions of your toolset
 - > check the manuals, the emitted warnings, ...

- > what do you **not** prove?
 - > the executed code does what you have in mind
 - > same mistake in the code and in the spec
 - > missed assumptions
 - > ... *we'll come back at the end*

- > Formal verification of a **JavaCard Virtual Machine** [Djoudi et al, 2021]
- > Common Criteria's **EAL7** certificate
- > Example of properties
 - > **header integrity**
 - > allocated object's header cannot be modified during a run
 - > **data integrity**
 - > allocated object's data can be modified only by the owner
 - > **data confidentiality**
 - > allocated object's data can be read only by the owner
- > generate $\approx 400,000$ ACSL annotations from ≈ 500 **MetACSL properties**
 - > all proved with Wp


 THALES

`https://julien-signoles.fr/teaching/tp-cyber/tp.md`

- > Frama-C in a VM
- > based on **Eva** (value analysis by abstract interpretation)
- > target: ciphering and deciphering functions of **Bacon code**
 - > from *Rosetta Code* project
 - > extended with a few (simple) vulnerabilities introduced
- > **technical goals:**
 - 1) **find all the vulnerabilities** with Eva et **fix them**
 - 2) **demonstrate there is no vulnerabilities anymore**, i.e., 0 alarm
 - > for a large class of input messages



- > **principles:**
 - > let **students work independently** as much as possible
 - > help them and provide explanation when necessary
- > **difficult** to apply this principle for a lab session on automatic analysis
 - > if the analysis is conclusive (no alarm or bug found), one has no clue why it concludes this way
 - > otherwise, hard to understand why it is not able to conclude
- > **solution:** step-by-step problem wording
 - > **preliminary exercise** on a simple example from the Juliet test suite
 - > **slow development**
 - > **many written explanations** about the analyzer's output
- > **not (yet) solvable by any LLM**
 - > the students must often describe then explain what they see on the Frama-C GUI

- 1 } Context
- 2 } Lecture and Lab Session
- 3 } Feedback

- > usually not a lot of feedback in French classes, yet a bit, e.g.:
“The students said again at the review meeting that your class was great!”

J.-F. Lalande, 15th April 2025 (private mail, translated from French)

- > live examples and demos make the class more engaging
 - > also help for time management, e.g. skipping a demo if late is easy
- > industrial applications demonstrate usefulness of formal methods
- > pointing out limitations makes the explanations more credible

- > they like working on a real open-source code
- > good level of difficulty
 - > not too easy, yet doable alone (if they take care to all the written advice)
 - > fine-tuning over several years: added questions and explanations when necessary
- > teacher's advice still needed from time to time
 - > point out what is written in the text
 - > explain some Frama-C outputs and/or behaviors
- > unintended positive side-effects: improve their C programming knowledge a bit
 - > computer arithmetics and integer overflows
 - > dynamic arrays and pointer arithmetics
 - > malloc may fail and return `null`
 - > error handling
- > some students “fix” code in the wrong way, leading to unremovable alarms
- > sparse VM issue (VM size, no ARM support, etc)

- > not as mature as the standard 3-hour version
- > shortening the lesson was easy
 - > same outline and global message
 - > mainly less demos and examples (which take time)
 - > yet message harder to deliver: less time to repeat the key insights regularly
- > shortening the lab session was more difficult
 - > remove preliminaries and move its relevant questions to the main exercise
 - > repetitive questions have also been removed
 - > now even a bit too easy for the best students
 - > also lack repetition for good learning
 - > critical to solve any VM issue very quickly

- > 3-hour lecture and 3-hour lab session about [Frama-C for cybersecurity](#)
- > **practical introduction** to formal methods and their applications to cybersecurity
- > **reusable material** for teachers familiar with Frama-C
 - > already successfully done once, without modifying the material
- > **adaptable to others formats**
 - > shorter lecture and shorter lab session
 - > 2-hour lecture in a doctoral summer school
 - > 1-hour seminar to master students

> Teachers who have invited me to teach Frama-C for cybersecurity:



O. Levillain
Telecom SP



J.-F. Lalande
CS Rennes



P. Wilke



S. Sadou
U. Bretagne Sud



G. Gogniat



O. Hermant
Mines Paris

> Original idea of the lab session:



V. Prevosto

Some examples:



N. Kosmatov